## Lecture 15: September 22

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We can view a partition of $n$ as a multiset $\left\{\lambda_{1}, \ldots, \lambda_{n}\right\}$ of elements choose from $\{1,2, \ldots\}$ so that $\lambda_{1},+\ldots+\lambda_{k}=n$ or better $w\left(\lambda_{1}\right)+\ldots w\left(\lambda_{k}\right)=n$. Multisets or unordered: this corresponds to the fact that we have a canonical way to write $\lambda$ namely in non-increasing order.
We can regard the parts as connected objects which there are exactly on for each weight (length) and a partition is then a multiset of connected objects.

### 15.1 Unlabeled Trees and Forests

Trees $=$ a connected acyclic graph on $n$ vertices. connected $\Rightarrow$ nonempty
Number of trees on $n$ vertices: $t_{n}=\sum_{k=1}^{\infty} t_{k} x^{k}$

$$
\begin{aligned}
& n=0, t_{n}=0 \\
& n=1, t_{n}=1 \\
& n=2, t_{n}=1 \\
& n=3, t_{n}=1 \\
& n=4, t_{n}=2
\end{aligned}
$$



$$
n=5, t_{n}=3
$$


$n=6, t_{n}=6$


A forest is a multiset of trees. If we let $f_{n}=$ the number of forests on $n$ vertices, then how can we compute $f_{n}$ ?
List all trees:



Give a list $\left(i_{1}, \ldots\right)$ of the number of copies of each tree. $i_{j}=$ number of copies of $T_{j}$. Then generating function then becomes

$$
\begin{aligned}
f(x) & =\sum_{n=0}^{\infty} f_{n} x^{n} \\
& =\left(1-x^{w\left(T_{1}\right)}\right)^{-1} \ldots\left(1-x^{w\left(T_{n}\right)}\right)^{-1} \\
& =\prod_{j=1}^{\infty}\left(1-x^{w\left(T_{j}\right)}\right)^{-1} .
\end{aligned}
$$

Grouping together by weight we see that we get

$$
f(x)=\prod_{m=1}^{\infty}\left(1-x^{m}\right)^{-t_{m}}
$$

We want to relate $f(x)$ and $t(x)$ :

$$
\begin{aligned}
f(x) & =\exp \left(-\sum_{m \geq 1} t_{m} \log \left(1-x^{m}\right)\right) \\
& =\exp \left(\sum_{m=1}^{\infty} \sum_{j=1}^{\infty} \frac{t_{m} x^{m j}}{j}\right) \\
& =\exp \left(\sum_{j=1}^{\infty} \frac{1}{j} \sum_{m=1}^{\infty} t_{m} x^{m j}\right) \\
& =\exp \left(\sum_{j=1}^{\infty} \frac{1}{j} \sum_{m=1}^{\infty} t_{m}\left(x^{j}\right)^{m}\right) \\
& =\exp \left(\sum_{j=1}^{\infty} \frac{t\left(x^{j}\right)}{j} .\right.
\end{aligned}
$$

Exactly the same method works and the same result holds true when $t(x)$ is the generating function for a set of connected objects and $f(x)$ is the generating function for multisets of connected objects. Example: Graphs on $n$ vertices, connected graphs on $n$ vertices.
Example: Rooted unlabelled trees, rooted forests.

$n=6, r_{n}=20$


3 ways,



Let $r(x)=\sum_{k \geq 1} r_{k} x^{k}$
$g(x)=\sum_{n \geq 0} g_{n} x^{n}, g_{n}=$ number of rooted forests on n vertices
$g(x)=\exp \left\{\sum \frac{1}{j} r\left(x^{j}\right)\right\}$
number of rooted forests on n vertices $=$ number of rooted trees on $\mathrm{n}+1$ vertices (bijection: chop off root) $\Rightarrow r(x)=x g(x)$
so $r(x)$ satisfies the functional equation $r(x)=x \exp \left\{\sum_{j=1}^{\infty} \frac{1}{j} r\left(x^{j}\right)\right\}$
Excercise: Use this functional equation to obtain enough information (a recurrence) to compute $r_{1}, r_{2}, \cdots, r_{10}$

Sets of connected objects:

$$
\begin{aligned}
h(x)=\sum_{n \geq 0} h_{n} x^{n} & =\prod_{j=1}^{\infty}\left(1+x^{j}\right)^{t_{j}} \\
& =\exp \left\{\sum_{m=1}^{\infty} t_{m} \log \left(1+x^{m}\right)\right\} \\
& =\exp \left\{\sum_{m=1}^{\infty} t_{m} \sum_{j=1}^{\infty}(-1)^{j+1} \frac{x^{m j}}{j}\right\} \\
& =\exp \left\{-\sum_{j=1}^{\infty} \frac{(-1)^{j}}{j} \sum_{m=1}^{\infty} t_{m} x^{m j}\right\} \\
& =\exp \left\{-\sum_{j=1}^{\infty} \frac{(-1)^{j}}{j} t\left(x^{j}\right)\right\} \\
& =\exp \left\{t(x)-\frac{1}{2}\left(x^{2}\right)+\frac{1}{3} t\left(x^{3}\right)-\cdots\right\}
\end{aligned}
$$

